

Exercises to the lecture
Semantics
Sheet 5

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Delivery until — at —

Exercise 5.1 (Böhm trees)

For the following λY terms over the signature

$$\Sigma = \{\text{if} : o \rightarrow o \rightarrow o \rightarrow o, \text{or}, +, - : o \rightarrow o \rightarrow o, \text{iszero}, \mathbf{a} : o \rightarrow o\},$$

create their Böhm trees up to depth 4, note their type and describe what they compute in the data domain \mathbb{N} . Interpret Y as the greatest fixed point.

- $Y\lambda x^o.x^o$
- $Y\lambda x^o.\mathbf{a}x^o$
- $Y\lambda x^o.+ x^o y^o$
- $Y\lambda F^{o \rightarrow o}\lambda a^o.$
 $\text{if} (\text{or} (\text{iszero} (- a^o 1)) (\text{iszero} a^o))$
 1
 $(+ (F^{o \rightarrow o} (- a^o 1) (F^{o \rightarrow o} (- a^o 2))))$

Exercise 5.2 (GFP-models for λY terms)

Proof the following Lemma from the lecture:

If $M =_{\beta\delta} N$, then for all GFP-models S and variable assignments ν , $\llbracket M \rrbracket_S^\nu = \llbracket N \rrbracket_S^\nu$.

Exercise 5.3 (TAC automata)

The definition for trees was $t : \{1, 2\}^* \rightarrow \Sigma \cup \{\Omega\}$, such that

1. If $uv \in \text{dom}(t)$ then $u \in \text{dom}(t)$
2. If $u \in \text{dom}(t)$ and $t(u) \in \Sigma_2$, then $u.1, u.2 \in \text{dom}(t)$
3. If $u \in \text{dom}(t)$ and $t(u) \in \Sigma_0 \cup \{\Omega\}$, then u is leaf.

A top-down tree automaton has the same syntactical definition as insightful TAC automata. However, its language is only defined for finite trees. When we interpret a given top-down tree automaton, we can find a pumping lemma. For this, we need to define for $u \in \text{dom}(t)$:

- $t_u :=$ the subtree rooted in u
- $t - u :=$ the tree t' , such that $\text{dom}(t') = \text{dom}(t) - u.\{1, 2\}^+$ and for each $u' \in \text{dom}(t')$, $t'(u') = t(u')$.
- $t +_u t'$, where $t(u) = t'(\varepsilon)$ is defined by $t(v)$ for $v \in \text{dom}(t - u)$ and by $t'(v)$ for $u.v$ with $v \in \text{dom}(t')$.

- if $t(\varepsilon) = t(u)$, $t^{n,u} = t +_u \underbrace{(t +_u (t +_u (\dots)))}_{n \text{ times}}$ and $t^{*,u} = \bigcup_{n \in \mathbb{N}} \{t^{n,u}\}$

Lemma. *Given a top-down tree automaton with n states and a tree $t \in L(A)$ with depth $> n$, there are vertices $u, u.v \in \text{dom}(t)$ such that $t = ((t - u) +_u ((t_u - v) +_v t_{uv}))$ and for each $n \in \mathbb{N}$, $((t - u) +_u ((t_u - v)^{n,v} +_{v^n} t_{uv})) \in L(A)$.*

Given a top-down tree automaton with n states accepting a language $L(A)$. What is the language of the same automaton interpreted as insightful TAC automaton?