

# Concurrency theory

## Exercise sheet 1

Roland Meyer, Elisabeth Neumann

TU Braunschweig  
Winter term 2017/18

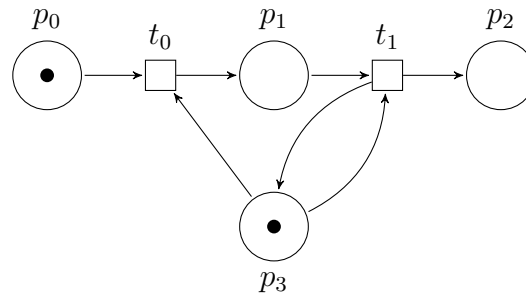
Out: November 29

Due: December 5

Submit your solutions until Wednesday, December 6, 12:00 am. You may submit in groups up to three persons.

### Exercise 1: The marking equation

Consider the following Petri net.



- Write down the connectivity matrix  $\mathbb{C}$  of the Petri net.
- Argue that the marking  $M_f = (0, 0, 1, 0)$  that has one token in  $p_3$  is not reachable from the initial marking  $M_0 = (1, 0, 0, 1)$ .
- Prove that the marking equation  $M_f - M_0 = \mathbb{C} \cdot c$  has a solution (i.e. there is a vector  $c \in \mathbb{N}^T$  satisfying the equation).

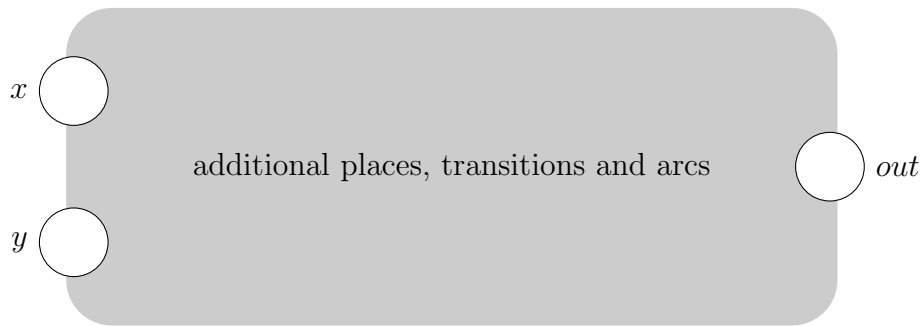
### Exercise 2: Petri net constructions

- Let  $(N, M_0, M_f)$  be a Petri net. Explain how to construct a Petri net  $(N', M'_0, M'_f)$  with  $M'_0(p) = 0$  for all places but a single place  $p'$  with  $M'_0(p') = 1$  and  $M'_f(p) = 0$  for all places such that  $M_f \in R(N, M_0)$  iff  $M'_f \in R(N', M'_0)$ .
- Let  $(N, M_0, M_f)$  be a Petri net. Explain how to construct a Petri net  $(N', M'_0, M'_f)$  such that  $M_f$  is coverable from  $M_0$  in  $N$  iff  $M'_f$  is reachable from  $M'_0$  in  $N'$ .
- Construct a Petri net  $N$  with only 3 places, a marking  $M_0$  and markings  $M_{c \wedge r}$ ,  $M_{\neg c \wedge \neg r}$  and  $M_{c \wedge \neg r}$  such that
  - $M_{c \wedge r}$  is reachable and coverable from  $M_0$ ,
  - $M_{\neg c \wedge \neg r}$  is neither reachable nor coverable, and
  - $M_{c \wedge \neg r}$  is coverable, but not reachable.

In each part of this exercise, argue briefly that your construction is correct.

### Exercise 3: Addition and multiplication

Consider the (incomplete) Petri net containing places  $x, y$  and  $out$  depicted below.



- a) Add places and transitions to the net such that any computation of the net starting in  $M_0(x) = m, M_0(y) = n, M_0(out) = 0$  terminates in a marking  $M_f$  with  $M_f(out) = m + n$ .  
(*Terminating* means that no transition is enabled anymore.)
- b) Add places and transitions to the net such that any computation of the net starting in  $M_0(x) = m, M_0(y) = n, M_0(out) = 0$  terminates in a marking  $M_f$  with  $M_f(out) \in \{0, \dots, m \cdot n\}$ .

In each part of this exercise, argue briefly that your construction is correct.