

Übungen zur Vorlesung
Concurrency Theory
Blatt 4

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Abgabe bis 18.06.2024 um 23:59 Uhr

Aufgabe 4.1 (Intersection non-emptiness of VASS reachability languages)

Let V_1, V_2 be two initialized VASS over Σ . The intersection non-emptiness problem asks whether $L(V_1) \cap L(V_2) \neq \emptyset$ holds. Show that intersection non-emptiness is decidable for VASS reachability languages.

Note: This is the first step in the procedure that decides regular separability.

Aufgabe 4.2 (VASS languages are closed under rational transductions)

Let V be an initialized VASS over Σ and $T \subseteq \Sigma^* \times \Gamma^*$ a rational transduction between Σ and Γ . Show that the language $T(L(V)) = \{T(w) \mid w \in L(V)\}$ is a VASS reachability language.

Aufgabe 4.3 (Büchi boxes)

Let $A = (Q, \Sigma, q_0, \delta, Q_F)$ be an NFA over an alphabet Σ . The transition equivalence relation $\sim_A \subseteq \Sigma^* \times \Sigma^*$ is defined by $v \sim_A w$ if for all $q, q' \in Q$ we have $q \xrightarrow{v} q'$ iff $q \xrightarrow{w} q'$.

- (a) Show that \sim_A is an equivalence relation with finitely many equivalence classes.
- (b) Show that for every word $w \in \Sigma^*$ the equivalence class $[w]_{\sim_A} := \{v \in \Sigma^* \mid v \sim_A w\}$ is a regular language.

Hint: First show that $L_{q,q'} := \{w \in \Sigma^ \mid q \xrightarrow{w} q'\}$ is regular.*

Aufgabe 4.4 (The Ackermann function)

The three-argument Ackermann function φ is defined recursively as follows.

$$\begin{aligned} \varphi: \mathbb{N}^3 &\rightarrow \mathbb{N} \\ \varphi(m, n, 0) &= m + n \\ \varphi(m, 0, 1) &= 0 \\ \varphi(m, 0, 2) &= 1 \\ \varphi(m, 0, x) &= m && \text{for } x > 2 \\ \varphi(m, n, x) &= \varphi(m, \varphi(m, n-1, x), x-1) && \text{for } n > 0 \text{ and } x > 0 \end{aligned}$$

- (a) Formally prove the following equalities using induction:

$$\varphi(m, n, 0) = m + n, \quad \varphi(m, n, 1) = m \cdot n, \quad \varphi(m, n, 2) = m^n.$$

(b) Nowadays, one usually considers the following two-parameter variant.

$$\begin{aligned} A: \mathbb{N}^2 &\rightarrow \mathbb{N} \\ A(0, n) &= n + 1 \\ A(m, 0) &= A(m - 1, 1) && \text{for } m > 0 \\ A(m, n) &= A(m - 1, A(m, n - 1)) && \text{for } m > 0 \text{ and } n > 0 \end{aligned}$$

For example, we have

$$A(1, 2) = A(0, A(1, 1)) = A(0, A(0, A(1, 0))) = A(0, A(0, A(0, 1))) = A(0, A(0, 2)) = A(0, 3) = 4.$$

Similar to this computation, write down a full evaluation of $A(2, 3)$.

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