

Übungen zur Vorlesung  
 Concurrency Theory  
 Blatt 3

Prof. Dr. Roland Meyer  
 Jan Grünke

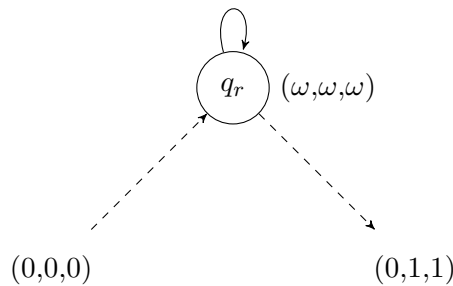
Abgabe bis 28.05.2024 um 23:59 Uhr

**Aufgabe 3.1** (Decomposition)

Decompose the following MGTS into a set of perfect MGTS.

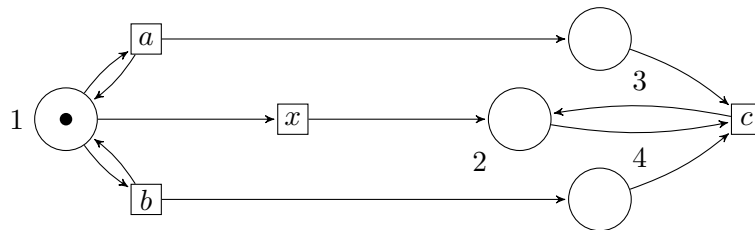
$$t_1 : (1,0,0) \quad t_3 : (0,-3,0)$$

$$t_2 : (1,5,-1) \quad t_4 : (-1,0,1)$$



**Aufgabe 3.2** (Rackoff's bound)

Consider the Petri net  $N = (\{1, 2, 3, 4\}, \{a, b, c, x\}, W)$  with multiplicities as depicted below. The initial marking of interest is  $M_0 = (1, 0, 0, 0)^T$  and the final marking is  $M_f = (1, 0, 10, 100)^T$ .



Compute the values  $m(3, M_0)$  and  $f(3)$  and argue why they are correct.

**Aufgabe 3.3** (Boundedness problem)

A Petri net is *bounded* if the set of reachable markings is finite, and *unbounded* otherwise. We consider the boundedness problem: Given a Petri net  $N = (S, T, W, M_0)$  with initial marking of size  $n$  decide whether  $N$  is bounded. Similar to the coverability problem the boundedness problem is decidable using coverability graphs. In this exercise we want

to show that it is decidable in EXPSPACE. Recall that a vector  $M \in \mathbb{Z}^S$  is  $i$ -bounded if the first  $i$  counters are positive and  $i$ - $r$ -bounded if the first  $i$  counters have values in  $\{1, \dots, r\}$ . We call a sequence of vectors  $M_1, \dots, M_k$  in  $\mathbb{Z}^S$  *self-covering* if there is some  $j \in [1, k]$  such that  $M_j < M_k$  (where we write  $M_1 < M_2$  if  $M_1 \leq M_2$  and  $M_1 \neq M_2$ ). A run is self-covering if the marking sequence of it is self-covering.

(a) Show that  $N$  is unbounded iff there is a self-covering run. More precisely, show that  $N$  is unbounded iff there are transition sequences  $\sigma_1, \sigma_2 \in T^*$  such that  $M_0[\sigma_1]M_1[\sigma_2]M_2$  and  $M_1 < M_2$ .

(b) For all  $r > 1$  and  $0 \leq i \leq |S|$  one can show the following: if there is a  $i$ - $r$ -bounded self-covering run, then there is a short  $i$ - $r$ -bounded self-covering run of length at most  $r^{n^c}$  for some constant  $c$ .

Use this to show  $g(i+1) \leq (2^n g(i))^{n^c}$  with

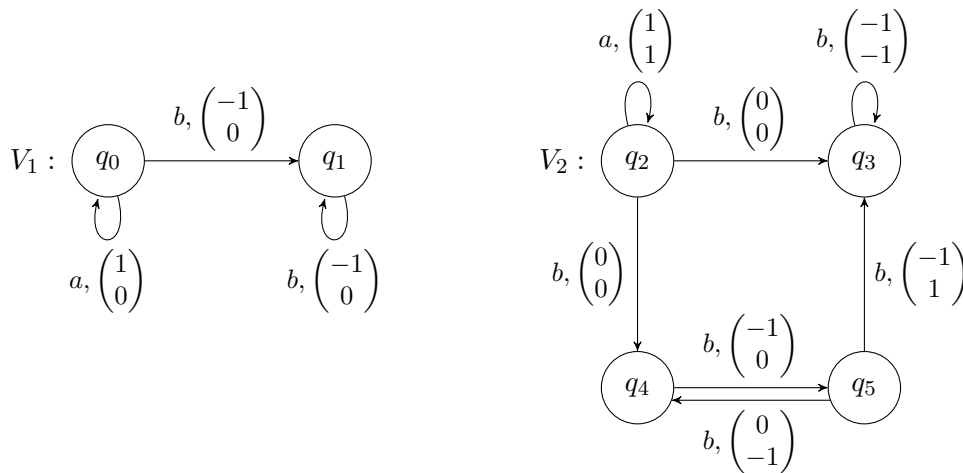
$$m(i, M) := \min \left\{ |\sigma| \mid \sigma \in \mathbb{Z}^{S^*} \text{ is } i\text{-bounded self-covering run from } M \right\}$$

$$g(i) := \max \left\{ m(i, M) \mid M \in \mathbb{Z}^S \right\}.$$

Give a bound for  $g(0)$ .

### Aufgabe 3.4 (Transducer Trick)

Let  $(V_1, (q_0, (0, 0)^T), (q_1, (2, 0)^T))$  and  $(V_2, (q_2, (1, 0)^T), (q_3, (1, 0)^T))$  be two initialized VASS over the alphabet  $\Sigma = \{a, b\}$  where  $V_1$  and  $V_2$  are given by:



Give another VASS  $V$  such that  $L(V)$  is regular separable from  $D_2$  (Dyck language with two counters) if and only if  $L(V_1)$  is regular separable from  $L(V_2)$ .

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