

8.1. Clearly  $\mathbb{Z}^2$  is partially ordered by  $R$  where

$sRs'$  iff.  $s \xrightarrow{w} s'$  for some  $w \in \{a, b\}^+$ .

Since each of  $T_1, T_2, T_3$  is a well-founded partial order over  $\mathbb{Z}^2$  and (by hypothesis)  $R \subseteq T_1 \cup T_2 \cup T_3$ , we can conclude from exercise 7.4 that  $R$  is well-founded.

This means that every infinite ordered sequence in  $R$  is stabilizing which proves the given program terminates.

8.2. (a)  $\Rightarrow$   $[u]_{\sim A} = [v]_{\sim A}$  means that for all  $q, q' \in Q$

$q \xrightarrow{u} q'$  iff.  $q \xrightarrow{v} q'$  and  $q \xrightarrow{u} \text{fin } q'$  iff.  $q \xrightarrow{v} \text{fin } q'$ .

This implies that  $R_{[u]_{\sim A}} = R_{[v]_{\sim A}}$  and  $R_{[u]_{\sim A}}^{\text{fin}} = R_{[v]_{\sim A}}^{\text{fin}}$ .

Therefore

$$\begin{aligned} \text{Box}(u) &= R_{[u]_{\sim A}} \cup R_{[u]_{\sim A}}^{\text{fin}} \\ &= R_{[v]_{\sim A}} \cup R_{[v]_{\sim A}}^{\text{fin}} = \text{Box}(v). \end{aligned}$$

$\Leftarrow$   $\text{Box}(u) = \text{Box}(v)$  means  $R_{[u]_{\sim A}} \cup R_{[u]_{\sim A}}^{\text{fin}} = R_{[v]_{\sim A}} \cup R_{[v]_{\sim A}}^{\text{fin}}$ ,

so  $R_{[u]_{\sim A}} = R_{[v]_{\sim A}}$  and  $R_{[u]_{\sim A}}^{\text{fin}} = R_{[v]_{\sim A}}^{\text{fin}}$ .

This implies that  $q \xrightarrow{u} q'$  iff.  $q \xrightarrow{v} q'$  and  $q \xrightarrow{u} \text{fin } q'$  iff.  $q \xrightarrow{v} \text{fin } q'$ , i.e.  $[u]_{\sim A} = [v]_{\sim A}$ .

(b)  $\Rightarrow$  If  $(q, q') \in \text{Box}(uv)$  it means  $\exists q''$  s.t.  $q \xrightarrow{u} q''$  and  $q'' \xrightarrow{v} q'$  so  $(q, q'') \in \text{Box}(u)$  and  $(q'', q') \in \text{Box}(v)$ .

Then, since  $R; S = \{(q, q') \in Q^2 \mid \exists q''. (q, q'') \in R \text{ and } (q'', q') \in S\}$  it means that  $\text{Box}(uv) \subseteq \text{Box}(u); \text{Box}(v)$ .

$\Leftarrow$  If  $(q, q') \in \text{Box}(u); \text{Box}(v)$  it means  $\exists q''$  s.t.  $(q, q'') \in \text{Box}(u)$  and  $(q'', q') \in \text{Box}(v)$ . This implies  $q \xrightarrow{u} q''$  and  $q'' \xrightarrow{v} q'$ .

Therefore  $q \xrightarrow{uv} q'$  and  $(q, q') \in \text{Box}(uv)$  which means that  $\text{Box}(u); \text{Box}(v) \subseteq \text{Box}(uv)$ .

8.2. (c) The following algorithm computes all  $\sim_A$  equivalence classes

```
resultSet = workSet = {Box( $\epsilon$ )}  $\cup \bigcup_{a \in \Sigma} \{Box(a)\}$ 
while workSet  $\neq \emptyset$  do
  pop Box( $v$ ) from workSet
  for each Box( $u$ ) in resultSet do
    if Box( $u$ ); Box( $v$ )  $\notin$  resultSet then
      add Box( $uv$ ) to resultSet and workSet
    end if
  end for
end while
return resultSet
```

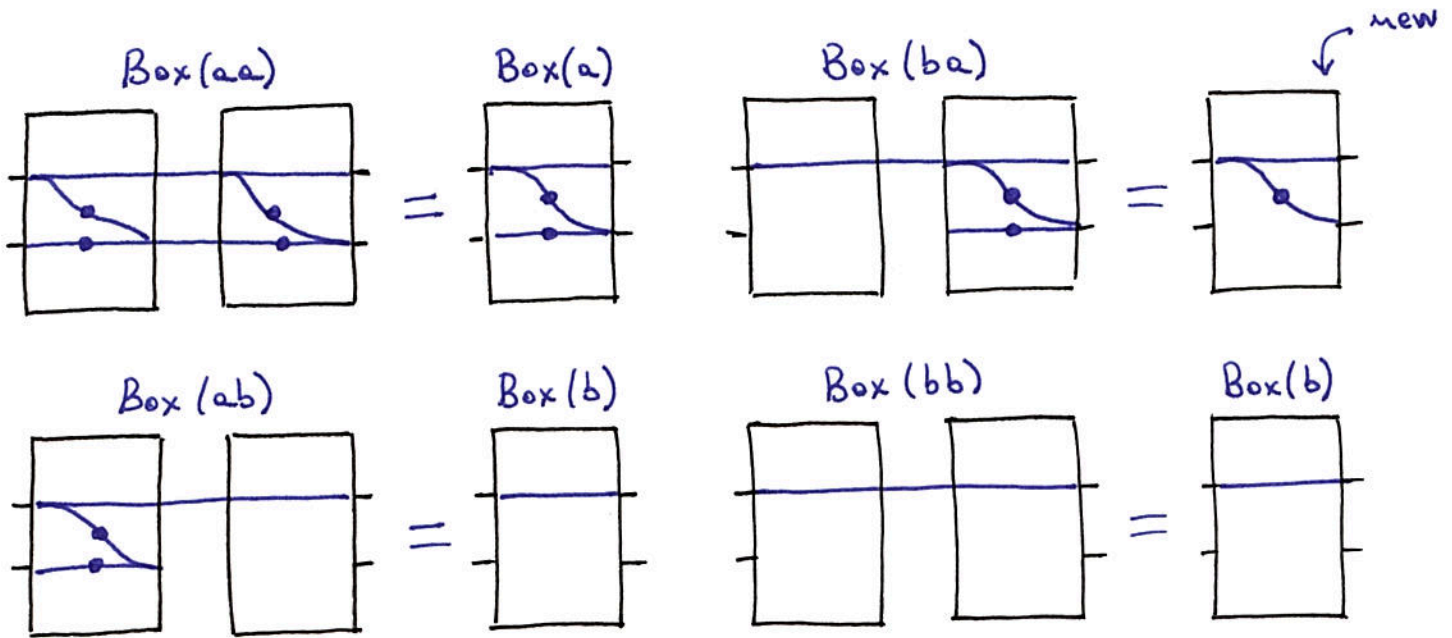
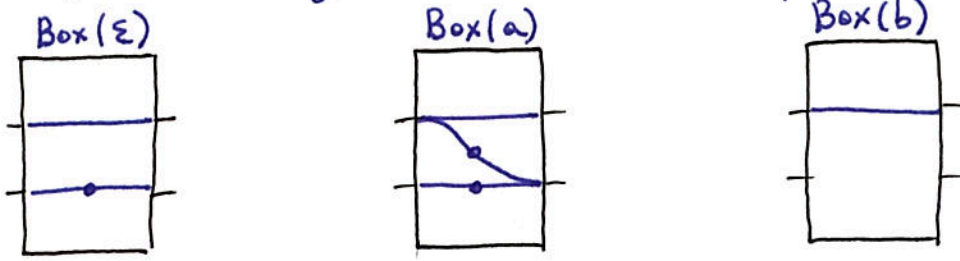
8.3. The following algorithm decides whether  $uv^w \in L(A)$

```
for each  $q \in Q_A$  do
  if  $q_0 \xrightarrow{uv} q$  and  $\exists n \in \mathbb{N}. q \xrightarrow{v^n} \text{fin } q$  then
    return true
  end if
end for
return false
```

Modifying the 8.2.(c) algorithm by adding the following will decide whether  $L(A) = \emptyset$ .

```
for each Box( $u$ ) and Box( $v$ ) in resultSet do
  if  $uv^w \in L(A)$  then
    return false
  end if
end for
return true
```

8.4. We first use algorithm 8.2.c to compute  $N_A$



$$\text{Box}(bab) = \text{Box}(bb) = \text{Box}(b)$$

$$\text{Box}(aba) = \text{Box}(ba)$$

$$\text{Box}(bba) = \text{Box}(ba)$$

$$\text{Box}(baba) = \text{Box}(bba) = \text{Box}(ba)$$

We then modify the algorithm for  $L(A) = \emptyset$  in 8.3 as follows:

$$\omega\text{-Set}_{L(A)} = \omega\text{-Set}_{\overline{L(A)}} = \emptyset$$

for each  $\text{Box}(u)$  and  $\text{Box}(v)$  in resultSet do

if  $uv^\omega \in L(A)$  then add  $([u]_{NA}, [v]_{NA})$  to  $\omega\text{-Set}_{L(A)}$

else add  $([u]_{NA}, [v]_{NA})$  to  $\omega\text{-Set}_{\overline{L(A)}}$

end if

end for

return  $\omega\text{-Set}_{L(A)}$  and  $\omega\text{-Set}_{\overline{L(A)}}$

By applying the algorithm we get  $L(A) = ([\varepsilon]_{NA} + [a]_{NA} + [b]_{NA} + [ba]_{NA})^\omega [a]_{NA}^\omega$ , respectively  $\overline{L(A)} = ([\varepsilon]_{NA} + [a]_{NA} + [b]_{NA} + [ba]_{NA}) ([b]_{NA}^\omega + [ba]_{NA}^\omega)$ .